

Electronic supplement for paper:

Tomas Skopal, David Hoksza:

*Improving the Performance of M-tree Family by Nearest-Neighbor Graphs*, AD-BIS 2007, Varna, Bulgaria, LNCS XXXX, Springer, 2007

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**Listing 1** (*k*-NN query algorithm)

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```
Node ChooseNode(PQueue PR) {
  let  $d_{min}(T(O_i^*)) = \min\{d_{min}(T(O_i))\}$ , considering all entries in PR
  remove entry  $[ptr(T(O_i^*)), d_{min}(T(O_i^*))]$  from PR
  return  $ptr(T(O_i^*))$ 
}

QueryResult kNNQuery(kNNQuery (Q, k), ordering heuristic  $\mathcal{H}$ ) {
  PR =  $\{[ptr(root), \infty]\}$ 
  for  $i = 1$  to  $k$  do
    NN[i] =  $[-, \infty]$  /*  $r_Q = NN[k].d_{max} = \infty$  */
    while PR is not empty do {
      NextNode = ChooseNode(PR)
      NodeSearch(NextNode, (Q, k),  $\mathcal{H}$ )
    }
  return NN
}
```

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The adaptation for M\*-tree is presented in Listing 2, where an M\*-tree node is processed. In addition to the original processing of M-tree node, the NN-graph filtering is inserted between the parent and basic filtering steps, however, we additionally store the already used sacrifices and use them repeatedly for filtering the same way as a new sacrifice is used (let us call this *recycled NN-graph filtering*). The rationale for this is specific for kNN processing – a sacrifice which was not successful in its first attempt could succeed after a sufficiently large decrease of the dynamically improving query radius  $r_Q$ .

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**Listing 2** (*k*-NN query algorithm, the NodeSearch)

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NodeSearch(Node  $N$ , kNNQuery ( $Q, k$ ), ordering heuristic  $\mathcal{H}$ ) {
  let  $P$  be the parent routing object of  $N$ 
    /* if  $N$  is root then  $\delta(O_i, P) = \delta(P, Q) = 0$  */
  let  $filtered$  be an array of boolean flags, size of  $filtered$  is  $|N|$ 
  set  $filtered[entry(O_i)] = false, \forall entry(O_i) \in N$ 
  let  $usedSacrifices = \emptyset$ 
  let  $SQ$  be a queue filled with all entries of  $N$ , ordered by  $\mathcal{H}(N)$ 
  if  $N$  is not a leaf then {
    while  $SQ$  not empty
      fetch  $rout(S_i)$  from the beginning of  $SQ$ 
      /* parent filtering */
      if  $|\delta(P, Q) - \delta(S_i, P)| > r_Q + r_{S_i}$  then
         $filtered[rout(S_i)] = true;$ 
      if not  $filtered[rout(S_i)]$  then {
        compute  $\delta(S_i, Q)$ 
        insert  $\langle S_i, \delta(S_i, Q) \rangle$  into  $usedSacrifices$ 
        if  $d_{min}(T(S_i)) \leq r_Q$  then {
          insert  $[ptr(T(S_i)), d_{min}(T(S_i))]$  to PR
          /* basic filtering */
          if  $d_{max}(T(S_i)) < r_Q$  then {
             $r_Q = \mathbf{NNUpdate}([- , d_{max}(T(S_i))])$ 
            remove PR requests for which  $d_{min}(T(S_i)) > r_Q$ 
          }
        }
      }
      /* NN-graph filtering */
       $NF = \emptyset$ 
      for each  $S_j$  in  $usedSacrifices$  do
         $NF = NF \cup \mathbf{FilterByNNGraph}(N, \langle S_j, \delta(S_j, Q) \rangle, (Q, r_Q))$ 
        move all entries in  $QS \cap NF$  to the beginning of  $QS$ 
    }
  } else { /*  $N$  is a leaf */
    while  $SQ$  not empty
      fetch  $grnd(S_i)$  from the beginning of  $SQ$ 
      /* parent filtering */
      if  $|\delta(P, Q) - \delta(S_i, P)| > r_Q$  then
         $filtered[grnd(S_i)] = true;$ 
      if not  $filtered[grnd(S_i)]$  then {
        compute  $\delta(S_i, Q)$ 
        insert  $\langle S_i, \delta(S_i, Q) \rangle$  into  $usedSacrifices$ 
        /* basic filtering */
        if  $\delta(S_i, Q) \leq r_Q$  then {
           $r_Q = \mathbf{NNUpdate}([- , d_{max}(T(S_i))])$ 
          remove PR requests for which  $d_{min}(T(S_i)) > r_Q$ 
        }
      }
      /* NN-graph filtering */
       $NF = \emptyset$ 
      for each  $\langle S_j, \delta(S_j, Q) \rangle$  in  $usedSacrifices$  do
         $NF = NF \cup \mathbf{FilterByNNGraph}(N, \langle S_j, \delta(S_j, Q) \rangle, (Q, r_Q))$ 
        move all entries in  $QS \cap NF$  to the beginning of  $QS$ 
    }
  }
}

```

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